

Towards Compact and Tractable Automaton-based Representations of Time Granularities

Ugo Dal Lago¹, Angelo Montanari², and Gabriele Puppis²

¹ Dipartimento di Scienze dell'Informazione, Università di Bologna
Mura Anteo Zamboni 7, 40127 Bologna, Italy
dallago@cs.unibo.it

² Dipartimento di Matematica e Informatica, Università di Udine
via delle Scienze 206, 33100 Udine, Italy
{montana,puppis}@dimi.uniud.it

Outline

- Time Granularities
- Representation Formalisms
- String-based and Automaton-based Approaches
- A Relevant Problem: The Granule Conversion Problem
- Optimizing Automaton-based Representations

Time Granularities

Motivations:

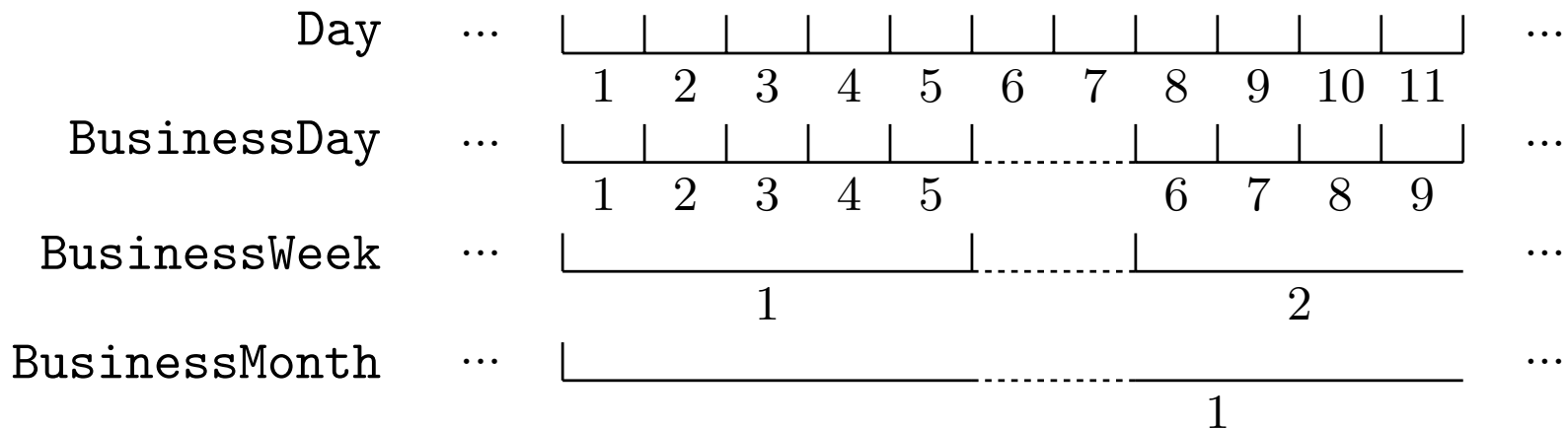
- **relational databases:**
to express temporal information at different time granularities, to relate different granules and convert associated data (queries)
- **artificial intelligence:**
to reason about temporal relationships, e.g, to check consistency and validity of temporal constraints at different time granularities (temporal CSPs)
- **data mining:**
to discover temporal relationships between collected events, to derive implicit information from such relationships

Time Granularities

Definition $G : \mathbb{Z} \rightarrow 2^T$ is a *granularity* iff

- $(T, <)$ is a linearly ordered set of temporal instants,
- $t_x < t_y$ whenever $x < y$, $t_x \in G(x)$, and $t_y \in G(y)$.

A *granule* of G is a non-empty set $G(x)$ and $x \in \mathbb{Z}$ is said to be its *label*.



Representation Formalisms (1)

We cannot finitely represent all the granularities over an infinite temporal domain, so we have to restrict ourselves to a proper subclass of structures.

Possible approaches to model time granularities:

- **algebraic one:**

relationships between granularities are represented by terms built up from a finite set of operators

(e.g., $\text{Week} = \text{Group}_7(\text{Day})$ in the Calendar Algebra, see Bettini et al. '00)

- **logical one:**

granularities are defined by models of formulas in a given language (e.g., PLTL, see Combi et al., '02)

Representation Formalisms (2)

- **string-based one:**
relationships between temporal instants and granules are encoded by sequences of symbols from a given alphabet (e.g., granspecs, see Wijsen '00)
- **automaton-based one:**
automata are exploited to encode string-based descriptions of time granularities (e.g., Single-string Automata, see Dal Lago '01)

We focus our attention on string-based and automaton-based approaches.

Fundamental Problems

- **Equivalence:**
the problem of establishing whether two different representations define the same granularity
- **Conversion:**
the problem of relating granules from different time granularities and converting associated data
- **Minimization:**
the problem of computing the smallest representation(s) for a given granularity
- **Optimization:**
the problem of computing the representation(s) on which crucial algorithms (e.g., conversion algorithms) run faster

String-based Approach (1)

Basic ingredients:

- a discrete temporal domain T
- restriction to *left bounded periodical* granularities, namely, granularities that have an initial granule and, ultimately, periodically group instants of the temporal domain
- a fixed alphabet $\{\blacksquare, \square, \}$, where
 - represents instants covered by some granule,
 - represents gaps within and between granules,
 - } separates granules and defines labels

String-based Approach (2)

Example. The infinite word

■ ■ ■ ■ ■ □ □ λ ■ ■ ■ ■ ■ □ □ λ ...

represents the granularity BusinessWeek in terms of Day.

Proposition. *Ultimately periodic words over $\{\blacksquare, \square, \lambda\}$ capture all the left bounded periodical granularities.*

Remark. Such strings can be finitely represented by pairs (*granspecs*) of prefixes and repeating patterns.

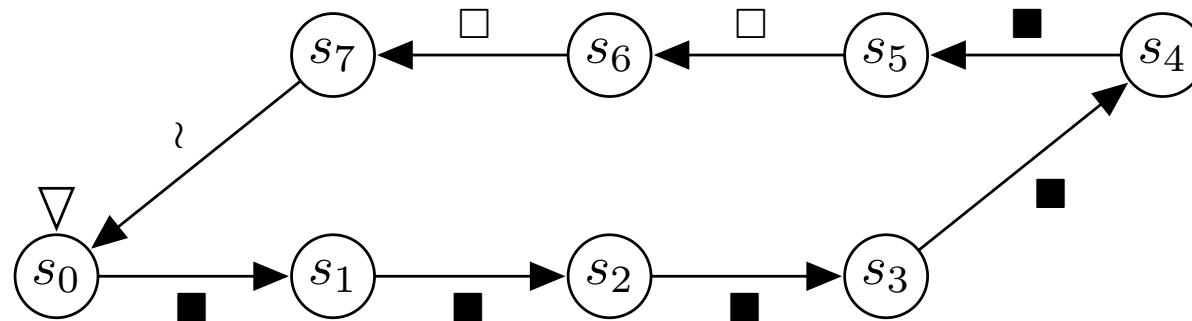
Example. $(\varepsilon, \blacksquare \blacksquare \blacksquare \blacksquare \blacksquare \square \square \lambda)$ is a granspec representing BusinessWeek in terms of Day.

Automaton-based Approach (1)

Connection between ultimately periodic words and automata:

Proposition. A single left bounded periodical granularity can be represented by an automaton recognizing a single string (SSA).

Example. An SSA representing BusinessWeek in terms of Day.



Problem. As for granspecs, a problem arises: such representations are *too large* with respect to inherently simple structure of granularities.

Automaton-based Approach (2)

Idea. We endow automata with *counters* and we use *primary and secondary transitions* to compactly encode redundancies of generated strings.

Remark. Suitable restrictions on the rules managing counter updates and transition activations guarantee

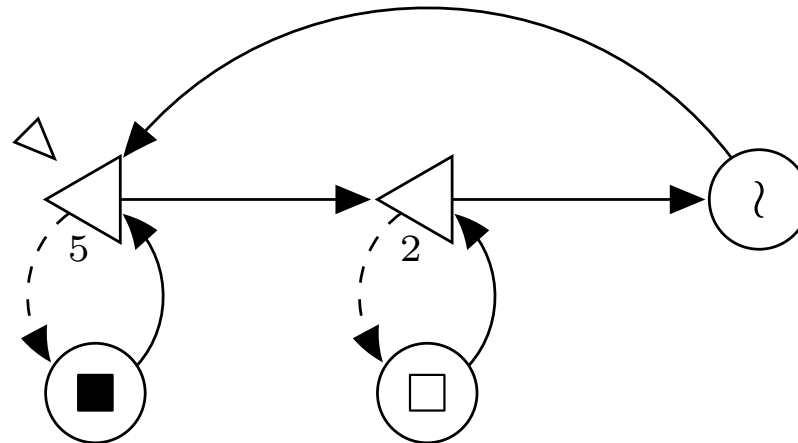
- the *same expressive power of SSA*;
- the *decidability of the equivalence problem*;
- *efficient manipulation* of representations (e.g. granule conversions).

Here the notion of *Restricted Labeled Automaton (RLA)* over finite and ultimately periodic words comes into play.

Automaton-based approach (3)

Distinctive features of RLA:

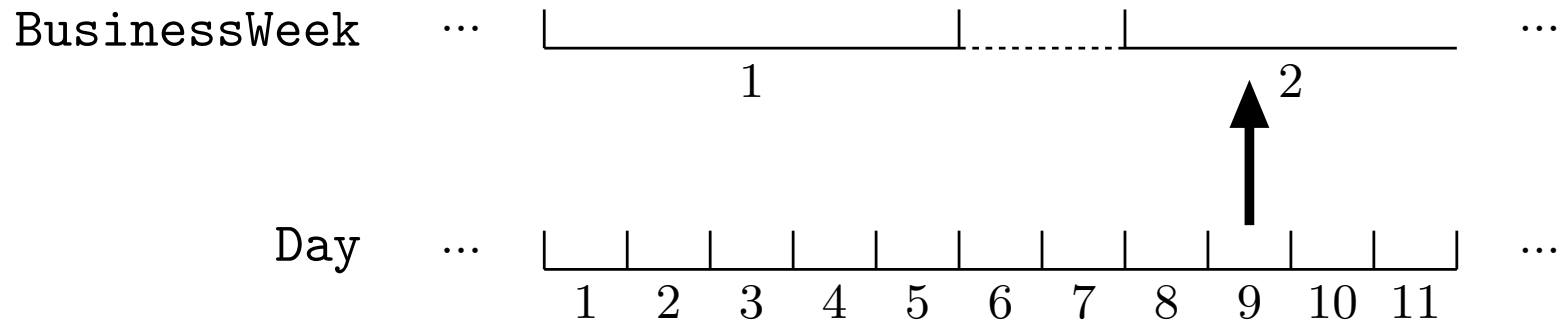
- states are partitioned into labeled and non-labeled states
- each non-labeled state is endowed with a single counter and it is the source of a secondary transition
- secondary transitions are activated iff the value of the corresponding counter is positive
- counters can only be decremented whenever they are positive, otherwise they are given their initial value



Granule Conversions (1)

Example (Granule Conversion Problem):

compute the label of the business week covering the 9-th day.



Granule conversions can be reduced to problems over strings/automata.

Example. In the given example, the solution is given by 1 plus the number of occurrences of λ in $(\blacksquare^5 \square^2 \lambda)^\omega$ until the 9-th occurrence of \blacksquare or \square (that is, 2).

Granule Conversions (2)

How RLA can be used to solve the problem?

- **straightforward solution:**
mimic the automaton transitions until the addressed occurrence has been reached
- **wiser solution:**
take advantage of nested repetitions in the run of the RLA in order to mimic *maximal periodic sequences* of transitions at once

⇒ The latter algorithm runs in *polynomial time* $\Theta(\|M\|)$, where M is the involved RLA and $\| \cdot \|$ is a suitable **complexity measure** envisaging the number of states and the structure of the transition functions.

The Optimization Problem

Problem. It is worth *minimizing the complexity* $\| \cdot \|$ in order to achieve the smallest running time for granule conversion algorithms.

Remark 1. There may be complexity-optimal automata which are not size-optimal, so

optimization problem \neq minimization problem.

Remark 2. There may be many different automata which are equivalent and complexity-optimal, so

there isn't a unique solution to the optimization problem.

Closure Properties of RLA (1)

Idea. We cope with the optimization problem by using *dynamic programming*.

Proposition. The class of RLA is closed under

- **Concatenation**

given two RLA recognizing u and v , it generates an RLA recognizing $u \cdot v$

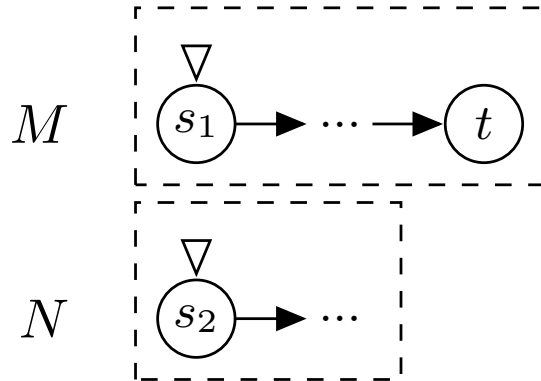
- **Iteration**

given an RLA recognizing u , it generates an RLA recognizing u^ω

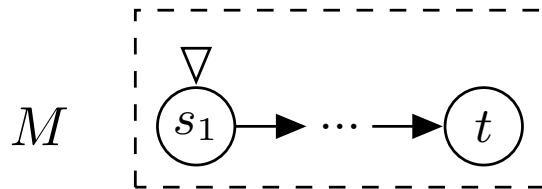
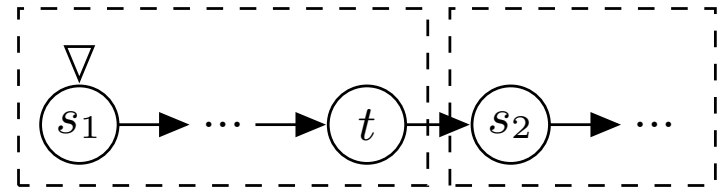
- **Repetition**

given an RLA recognizing u and a positive integer k , it generates an RLA recognizing u^k

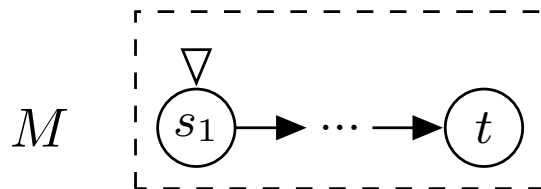
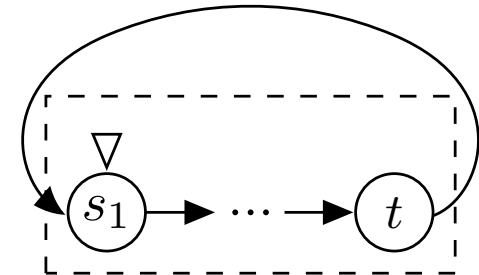
Closure Properties of RLA (2)



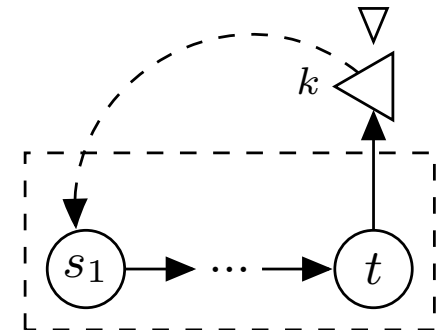
Concatenate(M, N)



Iterate(M)



Repeat(M, k)



Solving the Optimization Problem (1)

Remark. Unfortunately, there exist some RLA that cannot be generated by concatenation, iteration, and repetition.

However...

Proposition. For any RLA M , there exists an *equivalent* RLA M' , having the *same complexity*, which can be built starting from basic automata and using concatenation, iteration, and repetition.

⇒ We can restrict ourselves to a proper expressively complete subclass of compound automata, known as *Sharing Free RLA* (SFRLA for short).

Solving the Optimization Problem (2)

Theorem. Given a (finite or ultimately periodic) word u , there is a complexity-optimal SFRLA that recognizes u which is generated by combining smaller complexity-optimal SFRLA recognizing substrings of u .

Furthermore, we can restrict our search only to

a finite number of possible compositions

⇒ We can effectively build a complexity-optimal SFRLA recognizing a given string in a bottom-up fashion .

Such an optimization algorithm takes *polynomial time*.

Conclusions

We defined time granularities and we introduced an alternative automaton-based formalism for representing and reasoning about them.

We identified some crucial problems involving time granularities, with a special attention to the granule conversion problem.

We focused on the efficiency of conversion algorithms and we tackled the optimization of automaton-based representations.

Finally, we exploited non-trivial properties of RLA in order to solve the optimization problem.

It is an open question if a similar approach can be applied to solve the minimation problem.

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